

Cryptographic Challenges in and around Tor

Nick Mathewson
The Tor Project
9 January 2013

Summary

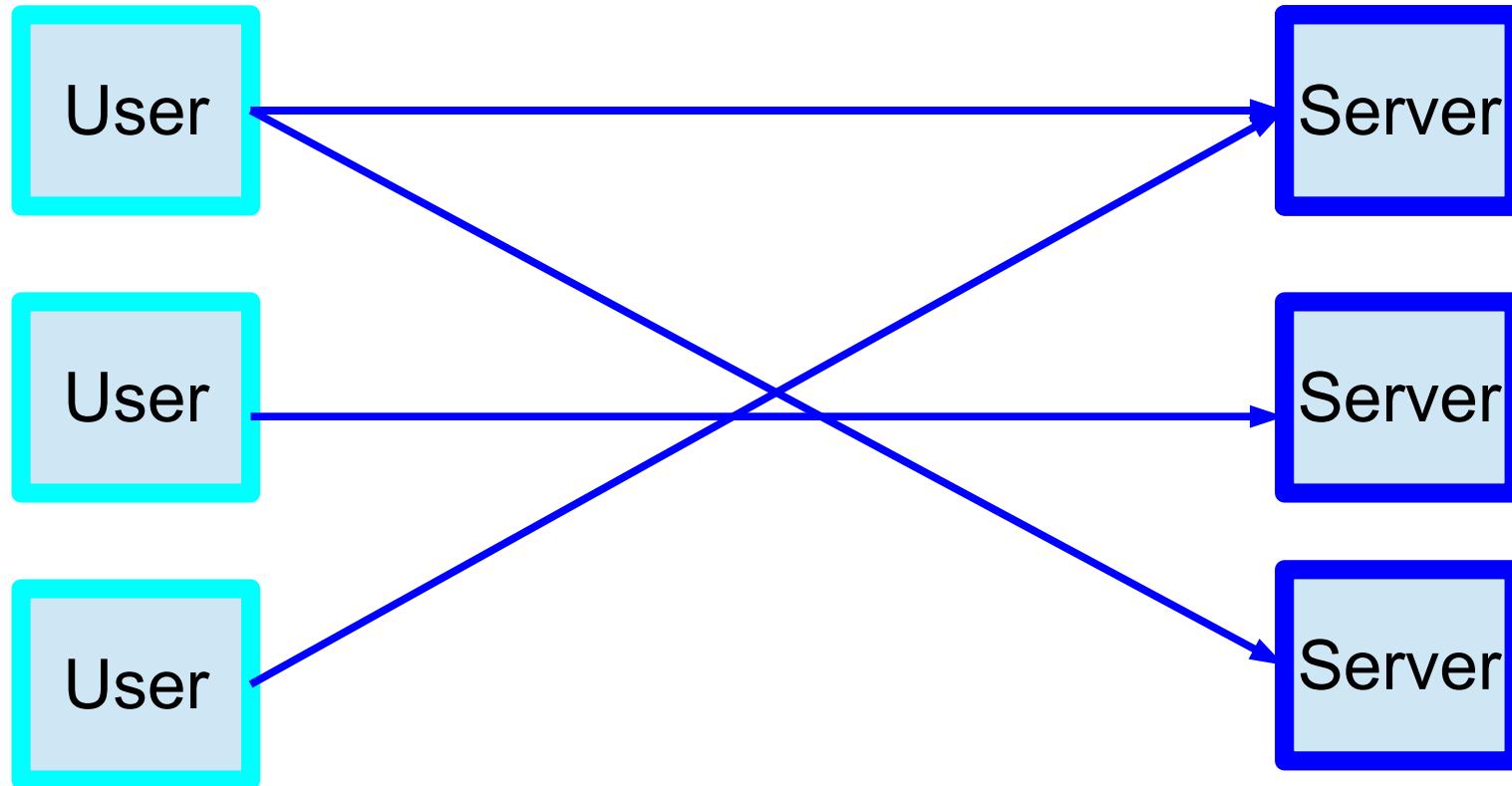
- Very quick Tor overview
- Tor's cryptography, and how it's evolving
- Various opportunities for more Tor crypto work

Disclaimer:

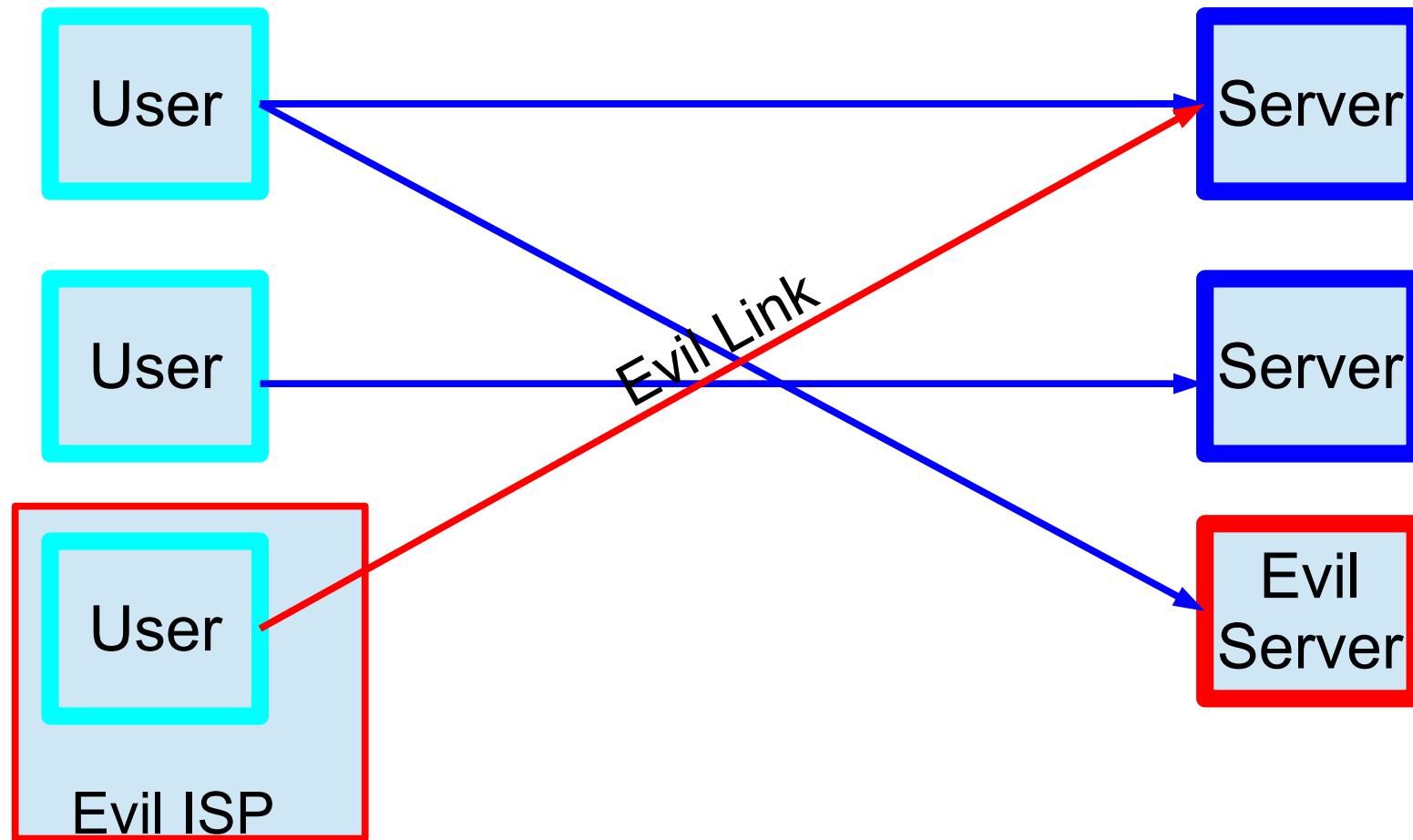
This is not exhaustive; these are only our most interesting crypto needs, not all of them; these are not our most urgent needs in general.

Part 1: Tor overview

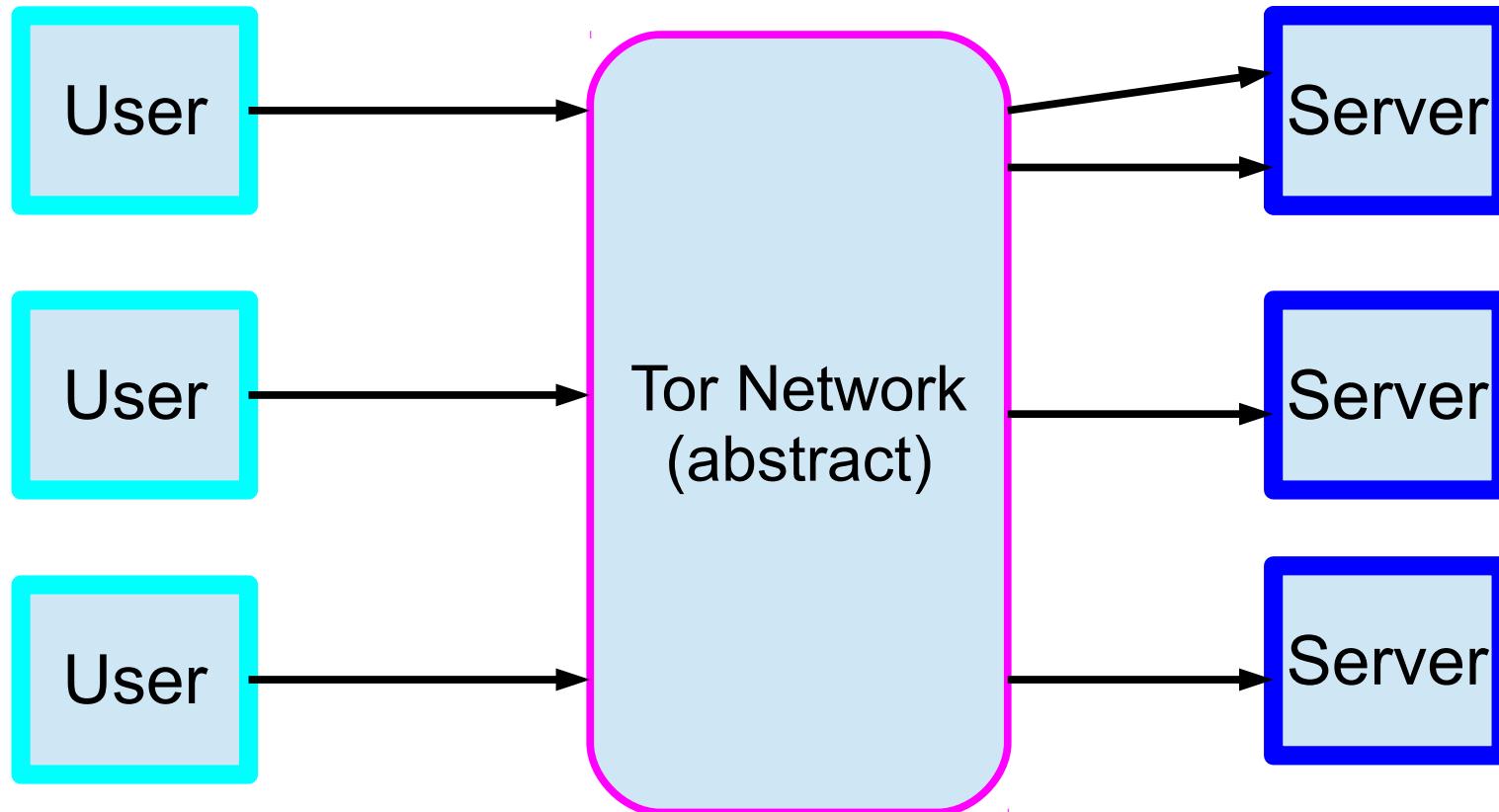
Ordinarily, traffic analysis and censorship are easy.



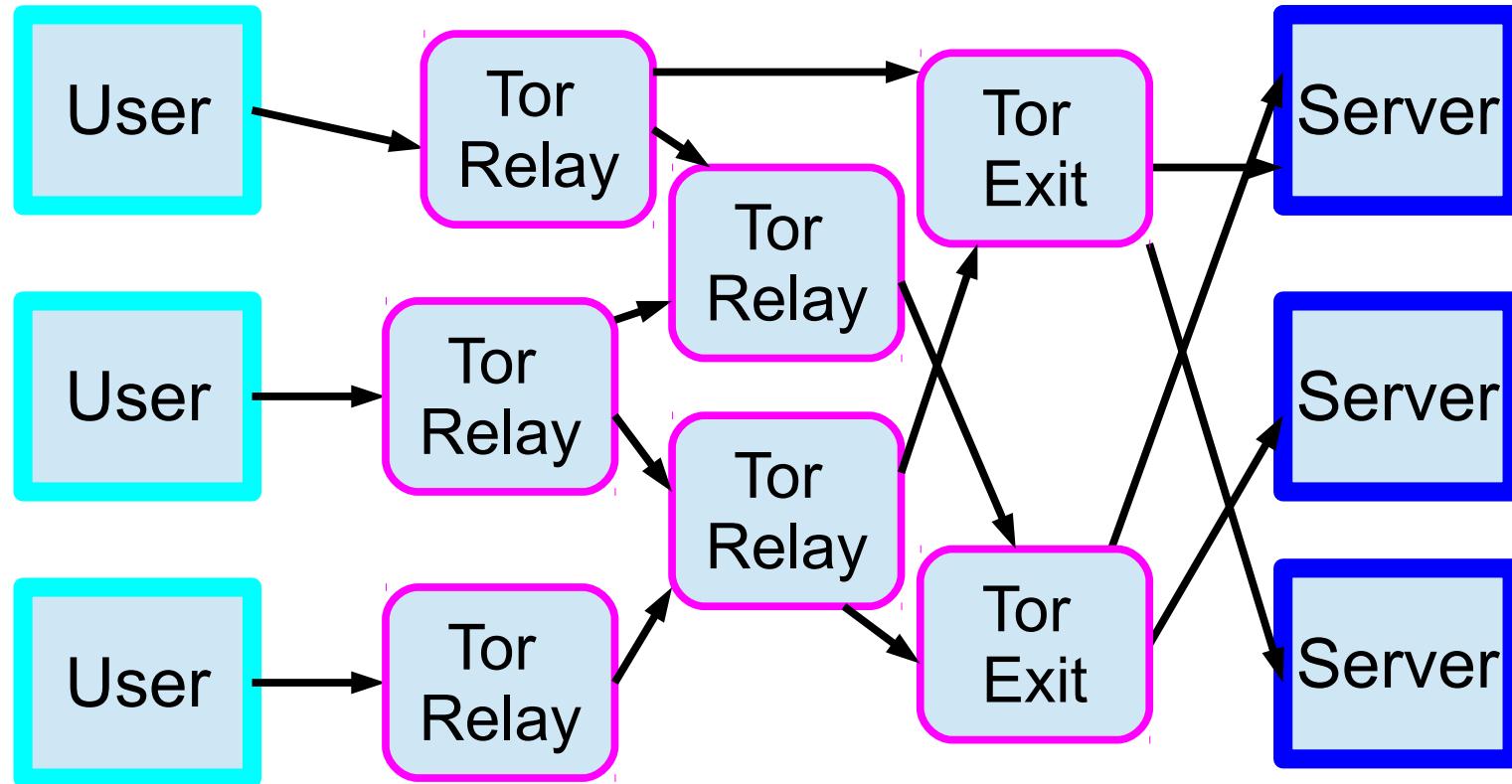
Ordinarily, traffic analysis and censorship are easy.



Tor makes traffic analysis and censorship harder...

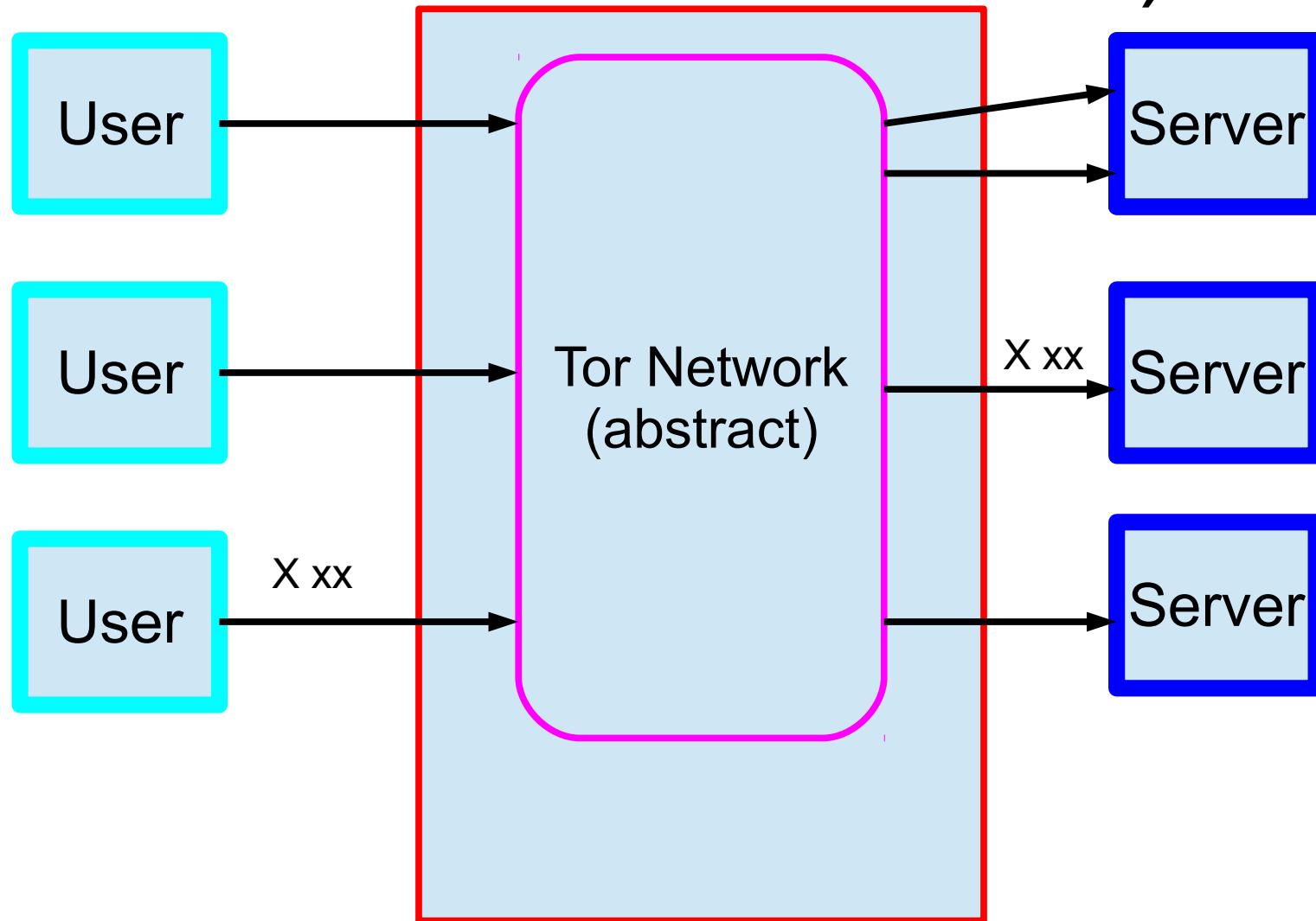


...by using a network of relays to anonymize traffic.



(Use non-public entry relays to resist censorship.)

(But an end-to-end traffic correlation attack still works.)

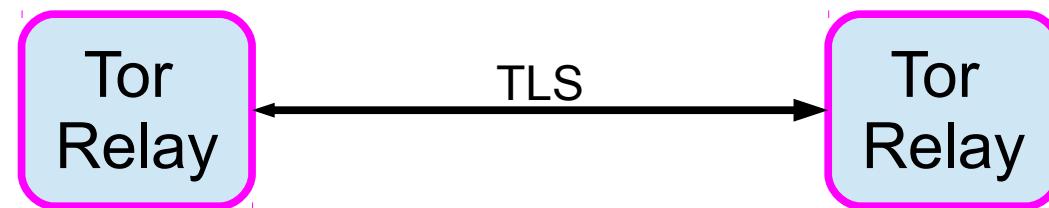
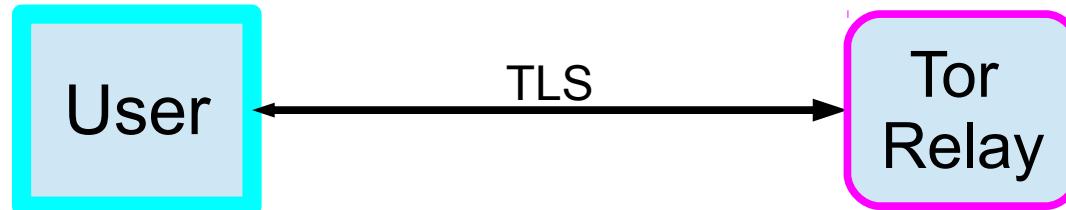


Tor is the largest deployed network of its kind

- 3000 relays
- 1000 public bridges
- > 2 GiB/sec
- > 500,000 users each day (estimated)
 - (With a pretty broad diversity of interest)

Part 2: Tor could use better crypto

Tor uses TLS for its link protocol...

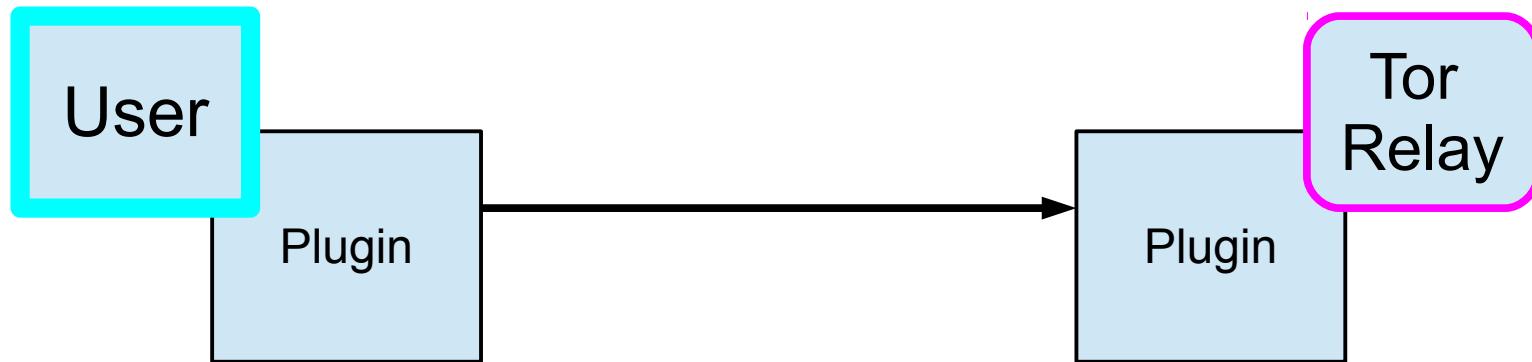


... with all the problems that entails.

- Easy to detect TLS variants based on:
 - Cipher choice
 - Certificate structure
 - List of extensions
- More secure: less common. Can't use any unpopular TLS feature.

(Did you know I have an effective veto over any new TLS features?)

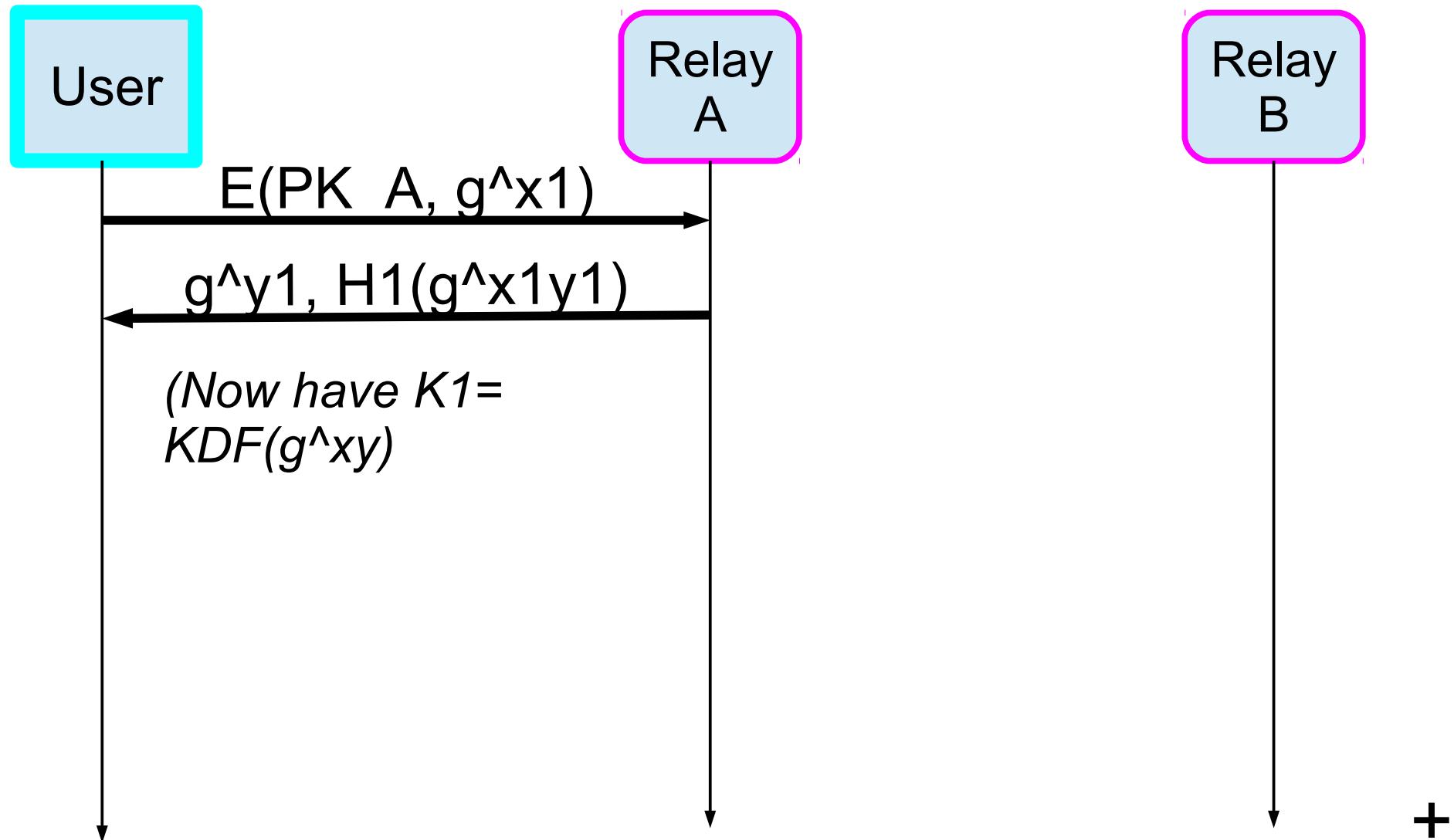
Maybe other link protocols are better
for anticensorship?



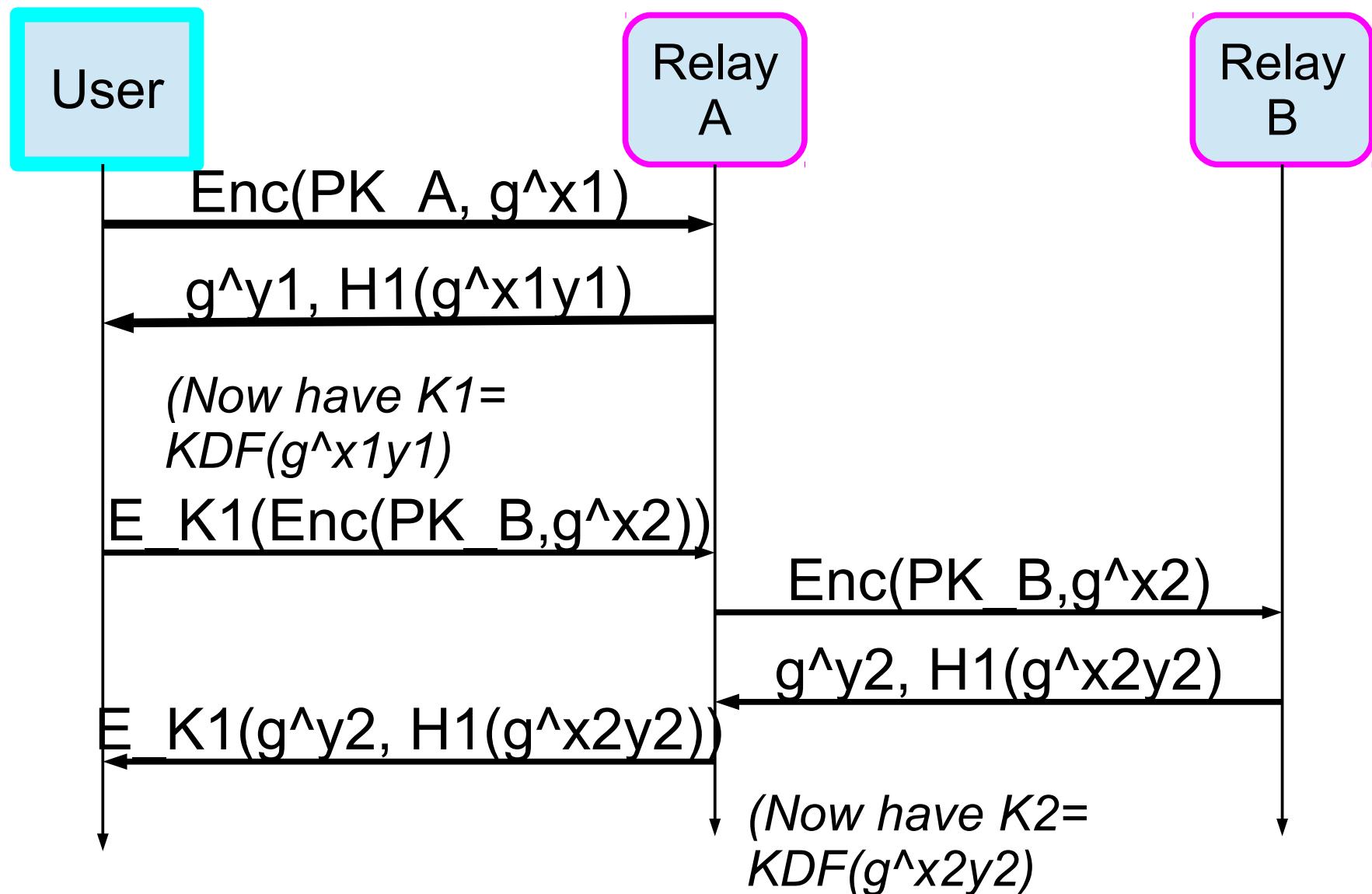
There are a number of these “Pluggable Transports” in development, but we need even more. *Even weak stego can help.*

...Do we still need “normal-looking” TLS?

Tor needs a one-way-authenticated handshake to build circuits



Tor needs a one-way-authenticated key exchange to build circuits



We're replacing this protocol...

- Original protocol (“TAP”) did hybrid encryption with RSA,DH-1024, badly. [Goldberg 2006]

- Replacement (“ntor”) does ***approximately***

C->S: g^x

S->C: $g^y, H1(\text{inp} = H(g^x g^y g^{xy} \dots))$

$K = \text{KDF}(H2(\text{inp}))$

[Goldberg, Stebila, Ustaoglu 2011]

(We're using DJB's curve25519 for DH group)

...and might replace it again

- Alternative (“ace”) does approximately:

$C \rightarrow S: g^{\wedge}x_1, g^{\wedge}x_2$

$S \rightarrow C: g^{\wedge}y$

$K = \text{KDF}(g^{\wedge}[bx_1 + yx_2])$

[Backes, Kate, Mohammedi 2012]

- Best choices will depend on implementation tweaks.
- Can you do better?

We should replace our old relay cell protocol...

- Used for symmetric crypto once we have shared keys.

Zeros (2)	Bad “MAC” (4)	Payload (503)
-----------	---------------	---------------

We should replace our old relay cell protocol...

- Used for symmetric crypto once we have shared keys.

Zeros (2)	Bad “MAC” (4)	Payload
		AES_CTR(Key1)
		AES_CTR(Key2)
		AES_CTR(Key3)

We should replace our old relay cell protocol...

- Used for symmetric crypto once we have shared keys.

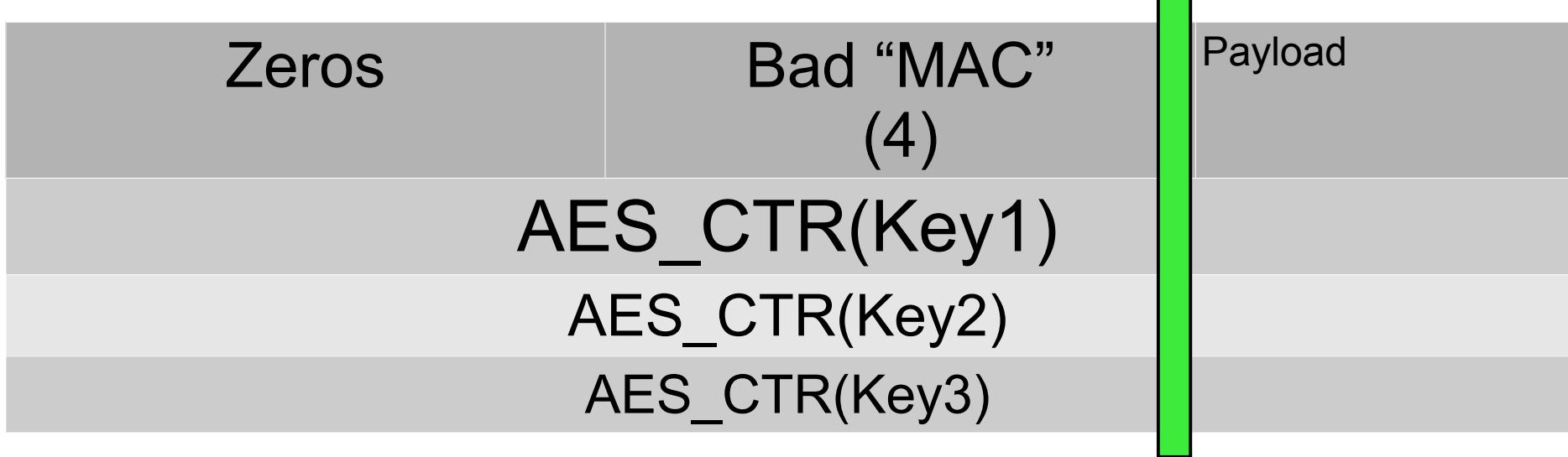
Zeros (2)	Bad “MAC” (4)	Payload
	AES_CTR(Key1)	
	AES_CTR(Key2)	
	AES_CTR(Key3)	

To handle a cell:

- Remove a layer of encryption.
- If Zeros == 0, and “MAC” = $H(Key3_M, Previous cells | Payload)$:
This cells is for us!
- Else, relay the cell

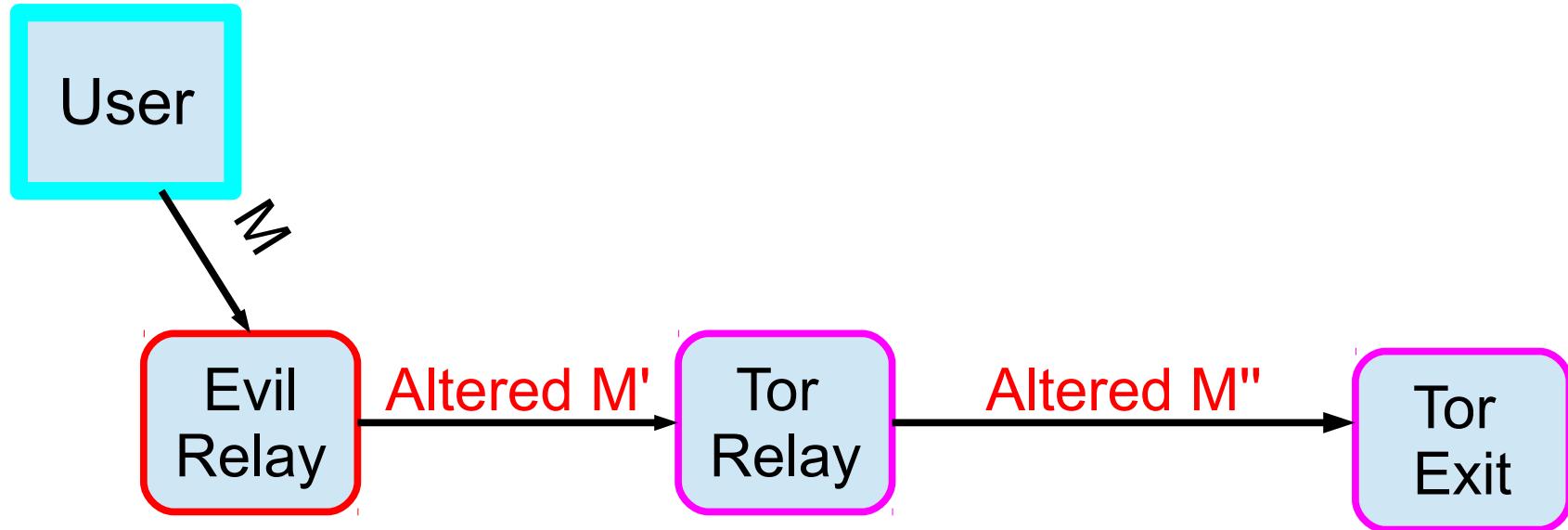
We should replace our old relay cell protocol...

- Used for symmetric crypto once we have shared keys.



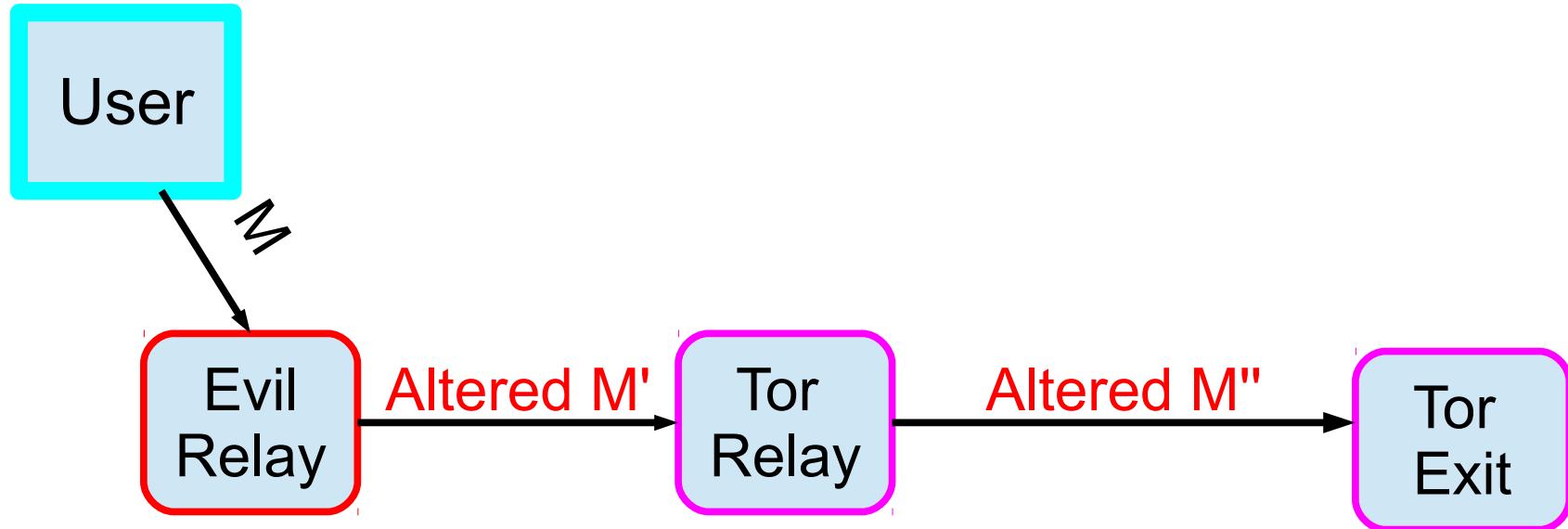
But this is malleable!

Hang on, does it matter that it's malleable?



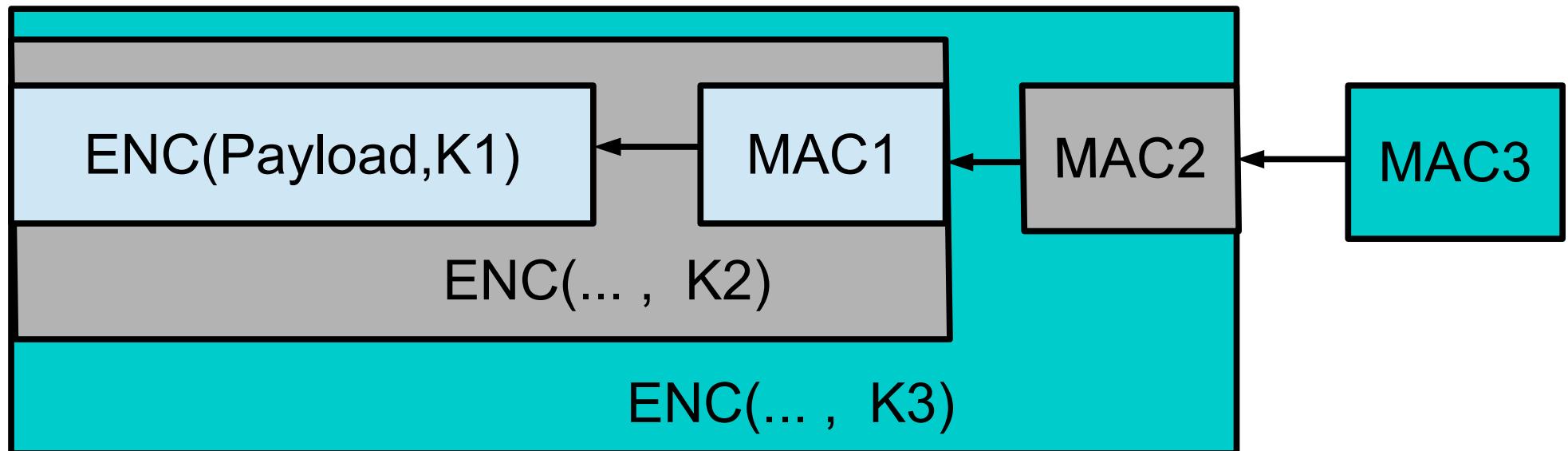
- Honest exit (probably) rejects M''
- Evil exit detects tag, but could just as easily do traffic correlation, for same result at less risk of detection.
- So, don't worry? (Dingledine, Mathewson, Syverson 2004)

Hang on, does it matter that it's malleable?

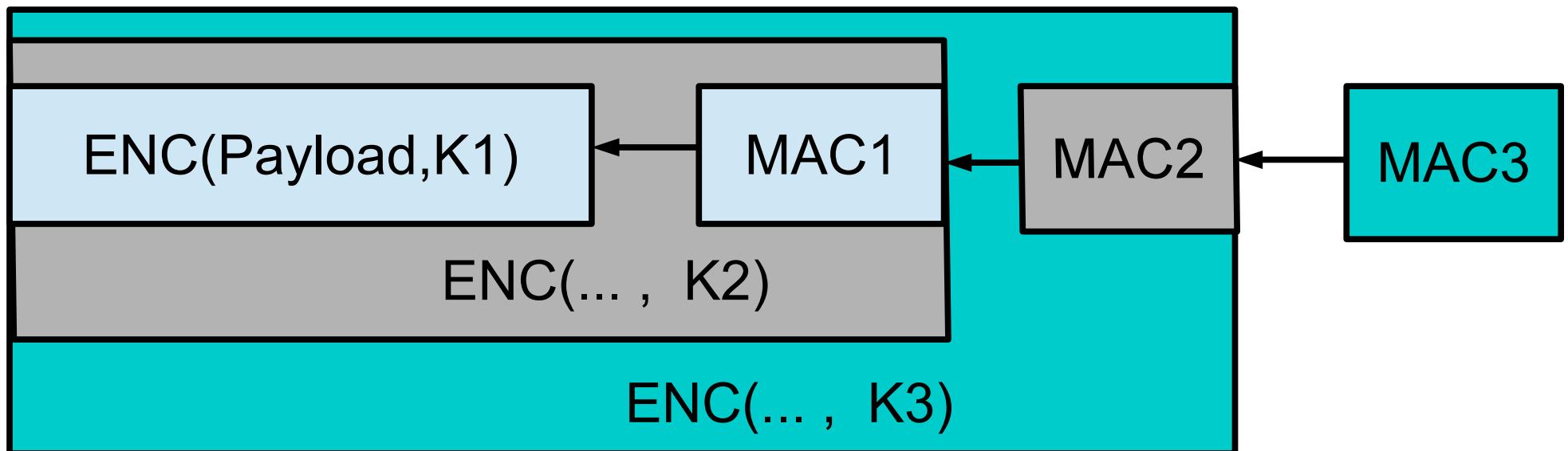


- Honest exit (probably) rejects M''
- Evil exit detects tag, but could ~~just as easily~~ do traffic correlation, ~~for the same result~~ at less risk of detection.
- *Actually, it's not so clear-cut.*

We could use an encrypt-and-mac structure



We could use an encrypt-and-mac structure

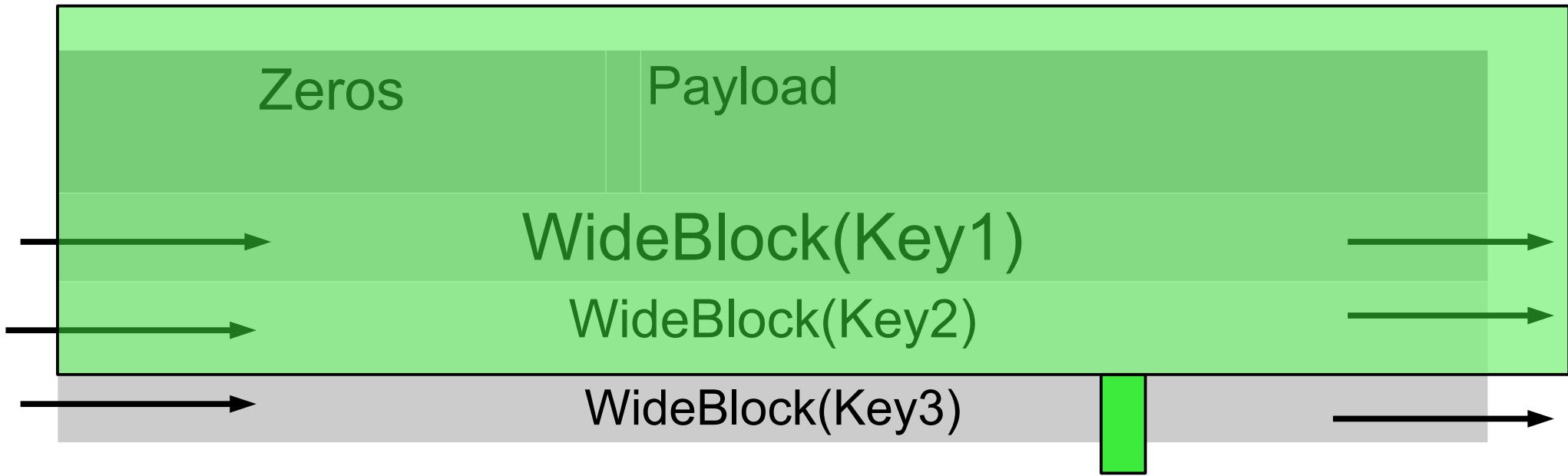


But that requires one MAC per hop, and leaks path length.

A chained wide-block cipher seems like a much better idea!



A chained wide-block cipher seems like a much better idea!



Any attempt to change the block renders the whole circuit unrecoverable.

What wide-block cipher to use?

- Not enough time to discuss all of them (LIONESS, CMC, XCB, HCTR, XTS, XEX, HCH, TET)
- Needs to be fast, proven, secure, easy-to-implement, non-patent-encumbered, side-channel-free,...
- One promising approach in progress by Bernstein, Sarkar, and Nandi – HFFH Feistel structure, fast, not yet finished.
- Other ideas?

Tor gets blocked too much.

- Some services mistake Tor for abuse
- Some services use IP blocking as a proxy for people-blocking, and can't *not* block Tor.
(Wikipedia edits, some IRC nets.)

Can we do better?

Provide a way for users to make themselves blockable.

- Slightly expensive pseudonyms?
 - (Expensive how? SA model?)
- Anonymous blacklistable credentials?
(Nymble, BNymble, BLACR, VERBS, Jack...)
 - Time to try this out in the wild?
 - What will we learn about their usability? Are they right?

There are more crypto issues in Tor

- Directory protocol
- Hidden service protocol
- Better DOS resistance
- SHA1, RSA1024 for node identity

Questions?

- See <https://www.torproject.org/> for links to documentation, specifications, and more info about various Tor issues.
- See <http://freehaven.net/anonbib/> for an incomplete but nonetheless useful anonymity bibliography.
- Grab me during a break for non-crypto Tor questions